

Cortex™-A9 NEON™ Media Processing Engine

Revision: r2p2

Technical Reference Manual

ARM®

Cortex-A9 NEON Media Processing Engine

Technical Reference Manual

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Release Information

The following changes have been made to this book.

Change history			
Date	Issue	Confidentiality	Change
04 April 2008	A	Non-Confidential	First release for r0p0
10 July 2008	B	Non-Confidential Restricted Access	Second release for r0p0
12 Dec 2008	C	Non-Confidential Restricted Access	First release for r1p0
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Product Status

The information in this document is final, that is for a developed product.

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Preface

This preface introduces the *Cortex-A9 NEON™ Media Processing Engine (MPE) Technical Reference Manual*. It contains the following sections:

- *About this book* on page xii
- *Feedback* on page xv.

About this book

This book is for the Cortex-A9 NEON MPE.

Product revision status

The *rn* identifier indicates the revision status of the product described in this book, where:

rn Identifies the major revision of the product.

pn Identifies the minor revision or modification status of the product.

Intended audience

This book is written for system designers, system integrators, and verification engineers who are designing a *System-on-Chip* (SoC) device that uses the Cortex-A9 NEON MPE. The book describes the external functionality of the Cortex-A9 MPE.

Using this book

This book is organized into the following chapters:

Chapter 1 *Introduction*

Read this for a high-level view of the Cortex-A9 NEON MPE and a description of its features.

Chapter 2 *Programmers Model*

Read this for a description of the Cortex-A9 NEON MPE system registers.

Chapter 3 *Instruction Timing*

Read this for a description of the cycle timings of instructions on the Cortex-A9 NEON MPE.

Appendix A *Revisions*

Read this for a description of the technical changes between released issues of this book.

Glossary Read this for definitions of terms used in this book.

Conventions

Conventions that this book can use are described in:

- *Typographical* on page xiii

Typographical

The typographical conventions are:

<i>italic</i>	Highlights important notes, introduces special terminology, denotes internal cross-references, and citations.
bold	Highlights interface elements, such as menu names. Denotes signal names. Also used for terms in descriptive lists, where appropriate.
monospace	Denotes text that you can enter at the keyboard, such as commands, file and program names, and source code.
<u>monospace</u>	Denotes a permitted abbreviation for a command or option. You can enter the underlined text instead of the full command or option name.
<i>monospace italic</i>	Denotes arguments to monospace text where the argument is to be replaced by a specific value.
monospace bold	Denotes language keywords when used outside example code.
< and >	Enclose replaceable terms for assembler syntax where they appear in code or code fragments. For example: MRC p15, 0 <Rd>, <CRn>, <CRm>, <Opcode_2>

Additional reading

This section lists publications by ARM and by third parties.

See Infocenter, <http://infocenter.arm.com>, for access to ARM documentation.

ARM publications

This book contains information that is specific to this product. See the following documents for other relevant information:

- *Cortex-A9 Technical Reference Manual* (ARM DDI 0388)
- *Cortex-A9 MPCore Technical Reference Manual* (ARM DDI 0407)
- *Cortex-A9 Floating-Point Unit Technical Reference Manual* (ARM DDI 0408)
- *Cortex-A9 MBIST Controller Technical Reference Manual* (ARM DDI 0414)
- *Cortex-A9 Configuration and Sign-Off Guide* (ARM DII 0146)

- *CoreSight™ PTM™-A9 Technical Reference Manual* (ARM DDI 0401)
- *CoreSight PTM-A9 Configuration and Sign-Off Guide* (ARM DII 0161)
- *CoreSight PTM-A9 Integration Manual* (ARM DII 0162)
- *CoreSight Program Flow Trace Architecture Specification* (ARM IHI 0035)
- *AMBA® Level 2 Cache Controller (L2C-310) Technical Reference Manual* (ARM DDI 0246)
- *L220 Cache Controller Technical Reference Manual* (ARM DDI 0329)
- *AMBA AXI Protocol Specification* (ARM IHI 0022)
- *AMBA Specification* (ARM IHI 0011)
- *ARM Architecture Reference Manual, ARMv7-A and ARMv7-R edition* (ARM DDI 0406)
- *RealView™ Compilation Tools Developer Guide* (ARM DUI 0203)
- *RealView ICE and RealView Trace User Guide* (ARM DUI 0155)
- *Intelligent Energy Controller Technical Overview* (ARM DTO 0005).

Other publications

This section lists relevant documents published by third parties:

- *ANSI/IEEE Std 754-1985, IEEE Standard for Binary Floating-Point Arithmetic.*

Feedback

ARM welcomes feedback on this product and its documentation.

Feedback on this product

If you have any comments or suggestions about this product, contact your supplier and give:

- The product name.
- The product revision or version.
- An explanation with as much information as you can provide. Include symptoms if appropriate.

Feedback on content

If you have comments on content then send an e-mail to errata@arm.com. Give:

- the title
- the number, ARM DDI 0409F
- the page numbers to which your comments apply
- a concise explanation of your comments.

ARM also welcomes general suggestions for additions and improvements.

Chapter 1

Introduction

This chapter introduces the Cortex-A9 implementation of the ARM Advanced SIMD media processing architecture. It contains the following sections:

- *About the Cortex-A9 NEON MPE* on page 1-2
- *Applications* on page 1-4
- *Product revisions* on page 1-5.

1.1 About the Cortex-A9 NEON MPE

The Cortex-A9 NEON MPE extends the Cortex-A9 functionality to provide support for the ARM v7 *Advanced SIMD* and *Vector Floating-Point v3* (VFPv3) instruction sets. The Cortex-A9 NEON MPE supports all addressing modes and data-processing operations described in the *ARM Architecture Reference Manual*.

The Cortex-A9 NEON MPE features are:

- SIMD and scalar single-precision floating-point computation
- scalar double-precision floating-point computation
- SIMD and scalar half-precision floating-point conversion
- 8, 16, 32, and 64-bit signed and unsigned integer SIMD computation
- 8 or 16-bit polynomial computation for single-bit coefficients
- structured data load capabilities
- dual issue with Cortex-A9 processor ARM or Thumb instructions
- independent pipelines for VFPv3 and Advanced SIMD instructions
- large, shared register file, addressable as:
 - thirty-two 32-bit S (single) registers
 - thirty-two 64-bit D (double) registers
 - sixteen 128-bit Q (quad) registers

See the *ARM Architecture Reference Manual* for details of the information that the MPE can hold in the different register formats.

The Cortex-A9 NEON MPE provides high-performance SIMD vector operations for:

- unsigned and signed integers
- single bit coefficient polynomials
- single-precision floating-point values.

The operations include:

- addition and subtraction
- multiplication with optional accumulation
- maximum or minimum value driven lane selection operations
- inverse square-root approximation
- comprehensive data-structure load instructions, including register-bank-resident table lookup.

See the *ARM Architecture Reference Manual* for details of the Advanced SIMD instructions.

Note

The Advanced SIMD architecture extension, its associated implementations, and supporting software, are commonly referred to as NEON™ technology.

1.1.1 VFPv3 architecture hardware support

The Cortex-A9 NEON MPE hardware supports single and double-precision add, subtract, multiply, divide, multiply and accumulate, and square root operations as described in the ARM VFPv3 architecture. It provides conversions between 16-bit, 32-bit and 64-bit floating-point formats and ARM integer word formats, with special operations to perform conversions in round-towards-zero mode for high-level language support.

All instructions are available in both the ARM and Thumb instruction sets supported by the Cortex-A9 processor family.

The Cortex-A9 NEON MPE provides an optimized solution in performance, power, and area for embedded and media intensive applications.

ARMv7 deprecates the use of VFP vector mode. The Cortex-A9 NEON MPE hardware does not support VFP vector operations. In this manual, the term vector refers to Advanced SIMD integer, polynomial and single-precision vector operations. The Cortex-A9 NEON MPE provides high speed VFP operation without support code. However, if an application requires VFP vector operation, then it must use support code. See the *ARM Architecture Reference Manual* for information on VFP vector operation support.

Note

This manual gives information specific to the Cortex-A9 NEON MPE implementation of the ARM Advanced SIMD and VFPv3 extensions. See the *ARM Architecture Reference Manual* for full instruction set and usage details.

1.2 Applications

The Cortex-A9 NEON MPE provides mixed-data type SIMD and high-performance scalar floating-point computation suitable for a wide spectrum of applications such as:

- personal digital assistants and smartphones for graphics, voice compression and decompression, user interfaces, Java interpretation, and *Just In Time* (JIT) compilation
- games machines for intensive three-dimensional graphics, digital audio and in-game physics effects such as gravity
- printers and *MultiFunction Peripheral* (MFP) controllers for high-definition color rendering
- set-top boxes for high-end digital audio and digital video, and interactive three-dimensional user interfaces
- automotive applications for engine management, power train computation, and in-car entertainment and navigation.

1.3 Product revisions

This section describes the differences in functionality between product revisions:

r0p0 - r1p0 There are no functionality changes although you must use the Cortex-A9 revision r1p0 design with revision r1p0 NEON MPE.

r1p0 - r2p0 There are no functionality changes although you must use the Cortex-A9 revision r2p0 design with this revision r2p0 NEON MPE.

r2p0 - r2p1 There are no functionality changes although you must use the Cortex-A9 revision r2p1 design with this revision r2p1 NEON MPE.

r2p1- r2p2 There are no functionality changes.

Chapter 2

Programmers Model

This chapter describes the Cortex-A9 NEON MPE programmers model. It contains the following sections:

- *About this programmers model* on page 2-2
- *New Advanced SIMD and VFP features* on page 2-4
- *IEEE754 standard compliance* on page 2-6
- *Supported formats* on page 2-7
- *Advanced SIMD and VFP register access* on page 2-8
- *Register summary* on page 2-13
- *Register descriptions* on page 2-14.

2.1 About this programmers model

This section introduces the VFPv3 and Advanced SIMD implementation provided by the Cortex-A9 NEON MPE. In addition it provides information on initializing the Cortex-A9 NEON MPE ready for application code execution.

In addition to features provided in previous combined Advanced SIMD and VFP implementations, the Cortex-A9 NEON MPE provides:

- half-precision, 16-bit, floating-point value conversion
- support for emulation of VFPv3-D32 and VFPv3-D16.

See the *ARM Architecture Reference Manual* for more information.

2.1.1 Advanced SIMD and VFP feature identification registers

The Cortex-A9 NEON MPE implements the ARMv7 Advanced SIMD and VFP extensions.

Software can identify these extensions and the features they provide, using the feature identification registers. The extensions are in the coprocessor space for coprocessors CP10 and CP11. You can access the registers using the VMRS and VMSR instructions, for example:

```
VMRS <Rd>, FPSID ; Read Floating-Point System ID Register
VMRS <Rd>, MVFR1 ; Read Media and VFP Feature Register 1
VMSR FPSCR, <Rt> ; Write Floating-Point System Control Register
```

See *Advanced SIMD and VFP register access* on page 2-8 for a description of the registers.

In addition there are coprocessor access control registers. See *Non-secure Access Control Register* on page 2-10 and *Coprocessor Access Control Register* on page 2-8.

2.1.2 Enabling Advanced SIMD and floating-point support

From reset, both the Advanced SIMD and VFP extensions are disabled. Any attempt to execute either a NEON or VFP instruction results in an Undefined Instruction exception being taken. To enable software to access Advanced SIMD and VFP features ensure that:

- Access to CP10 and CP11 is enabled for the appropriate privilege level. See *Coprocessor Access Control Register* on page 2-8.
- If Non-secure access to the Advanced SIMD features or VFP features is required, the access flags for CP10 and CP11 in the NSACR must be set to 1. See *Non-secure Access Control Register* on page 2-10.

In addition, software must set the FPEXC.EN bit to 1 to enable most Advanced SIMD and VFP operations. See *Floating-Point Exception Register* on page 2-17.

When Advanced SIMD and VFP operation is disabled because FPEXC.EN is 0, all Advanced SIMD and VFP instructions are treated as undefined instructions except for execution of the following in privileged modes:

- a VMSR to the FPEXC or FPSID register
- a VMRS from the FPEXC, FPSID, MVFR0 or MVFR1 registers.

Example 2-1 shows how to enable the Advanced SIMD and VFP in ARM *Unified Assembly Language* (UAL). This code must be executed in privileged mode.

Example 2-1 Enabling Advanced SIMD and VFP

```

MRC p15,0,r0,c1,c0,2 ; Read CPACR into r0
ORR r0,r0,#(3<<20) ; OR in User and Privileged access for CP10
ORR r0,r0,#(3<<22) ; OR in User and Privileged access for CP11
BIC r0, r0, #(3<<30) ; Clear ASEDIS/D32DIS if set
MCR p15,0,r0,c1,c0,2 ; Store new access permissions into CPACR
ISB ; Ensure side-effect of CPACR is visible
MOV r0,#(1<<30) ; Create value with FPEXC (bit 30) set in r0
VMSR FPEXC,r0 ; Enable VFP and SIMD extensions

```

At this point the Cortex-A9 processor can execute Advanced SIMD and VFP instructions.

———— Note —————

Operation is Unpredictable if you configure the *Coprocessor Access Control Register* (CPACR) such that CP10 and CP11 do not have identical access permissions.

2.2 New Advanced SIMD and VFP features

The Cortex-A9 NEON MPE implements the following new features in the ARMv7 Advanced SIMD and VFP architectures:

- *Half-precision floating-point value conversion*
- *Independent Advanced SIMD and VFP disable*
- *Dynamically configurable VFP register bank size.*

Full details of the new instructions and control register fields are in the *ARM Architecture Reference Manual*.

2.2.1 Half-precision floating-point value conversion

The half-precision floating-point value conversion adds support for both IEEE and the common graphic representation, referred to as alternative-half-precision representation, of 16-bit floating-point values. This provides a smaller memory footprint for applications requiring large numbers of lower-precision floating-point values to be stored, while avoiding the overhead of conversion in software.

Additional VFP and Advanced SIMD instructions enable conversion of both individual values and vectors of values to and from single-precision floating-point representation. These values can then be processed using the rest of the VFP and Advanced SIMD instructions.

See *Floating-Point Status and Control Register* on page 2-15 for how to select IEEE or alternative half-precision modes.

2.2.2 Independent Advanced SIMD and VFP disable

The independent Advanced SIMD disables permit Cortex-A9 implementations with Cortex-A9 NEON MPE to behave as though only the VFP extension were present. This lets you enable optimal operating-system task scheduling between Cortex-A9 multi-processor clusters containing both Cortex-A9 NEON MPE and floating-point only units.

The Cortex-A9 processor provides support for preventing use of this feature through Non-secure access. See *Non-secure Access Control Register* on page 2-10 and *Coprocessor Access Control Register* on page 2-8.

2.2.3 Dynamically configurable VFP register bank size

The dynamically configurable VFP register bank size provides additional support for both VFPv3-D16 and VFPv3-D32 mixed multiprocessor clusters. Cortex-A9 NEON MPE implements thirty-two 64-bit double-precision registers. VFP-only

implementations are only required to support sixteen double-precision registers. This register bank disable control enables emulation of a 16-entry double-precision register file, providing both enhanced compatibility and more flexible task scheduling.

Additional control is provided in the Non-Secure Access Control Register. See *Non-secure Access Control Register* on page 2-10 and *Coprocessor Access Control Register* on page 2-8.

2.3 IEEE754 standard compliance

The IEEE754 standard provides a number of implementation choices. The *ARM Architecture Reference Manual* describes the choices that apply to the Advanced SIMD and VFPv3 architectures.

The Cortex-A9 NEON MPE implements the ARMv7 Advanced SIMD and VFP extensions. It does not provide hardware support for the following IEEE754 operations:

- remainder
- round floating-point number to nearest integer-valued in floating-point number
- binary-to-decimal conversion
- decimal-to-binary conversion
- direct comparison of single-precision and double-precision values
- any extended-precision operations.

2.4 Supported formats

Table 2-1 shows the formats supported for each of the Advanced SIMD and VFPv3 instruction sets implemented by the Cortex-A9 NEON MPE. All signed integers are two's complement representations.

Table 2-1 Supported number formats

Format	Advanced SIMD	VFPv3
8-bit signed/unsigned integer	Yes	No
16-bit signed/unsigned integer	Yes	No
32-bit signed/unsigned integer	Yes	Yes ^a
64-bit signed/unsigned integer	Yes	No
16-bit half-precision floating-point	Yes ^a	Yes ^a
32-bit single-precision floating-point	Yes	Yes
64-bit double-precision floating-point	No	Yes
8-bit polynomials	Yes	No
16-bit polynomials	Yes	No

a. For conversion purposes only.

2.5 Advanced SIMD and VFP register access

Table 2-2 shows the system control coprocessor registers, accessed through CP15, that determine access to Advanced SIMD and VFP registers, where:

- CRn is the register number within CP15
- Op1 is the Opcode_1 value for the register
- CRm is the operational register
- Op2 is the Opcode_2 value for the register.

Table 2-2 Coprocessor Access Control registers

CRn	Op1	CRm	Op2	Name	Description
c1	0	c0	2	CPACR	See <i>Coprocessor Access Control Register</i>
c1	0	c1	2	NSACR	See <i>Non-secure Access Control Register</i> on page 2-10

2.5.1 Coprocessor Access Control Register

The CPACR Register sets access rights for the coprocessors CP10 and CP11, that enable the Cortex-A9 NEON MPE functionality. This register also enables software to determine if a particular coprocessor exists in the system.

The CPACR Register is:

- a read/write register common to Secure and Non-secure states
- accessible in privileged modes only
- has a reset value of 0.

Figure 2-1 shows the CPACR Register bit assignments.

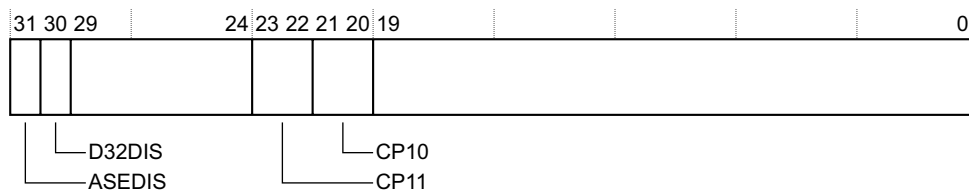


Figure 2-1 CPACR Register bit assignments

Table 2-3 shows the CPACR Register bit assignments.

Table 2-3 Coprocessor Access Control Register bit assignments

Bits	Field	Function
[31]	ASEDIS	Disable Advanced SIMD extension functionality: b1 = Disables all instruction encodings identified in the <i>ARM Architecture Reference Manual</i> as being part of the Advanced SIMD extensions but that are not VFPv3 instructions. b0 = No instructions are disabled.
[30]	D32DIS	Disable use of D16-D31 of the VFP register file: b1 = Disables all instruction encodings identified in the <i>ARM Architecture Reference Manual</i> as being VFPv3 instructions if they access any of registers D16-D31. b0 = No instructions are disabled.
[29:24]	-	See the <i>Cortex-A9 Technical Reference Manual</i> .
[23:22]	CP11	Defines access permissions for the coprocessor. Access denied is the reset condition and is the behavior for nonexistent coprocessors. b00 = Access denied. Attempted access generates an Undefined Instruction exception. b01 = Privileged mode access only. b10 = Reserved. b11 = Privileged and User mode access.
[21:20]	CP10	Defines access permissions for the coprocessor. Access denied is the reset condition and is the behavior for nonexistent coprocessors. b00 = Access denied. Attempted access generates an Undefined Instruction exception. b01 = Privileged mode access only. b10 = Reserved. b11 = Privileged and User mode access.
[19:0]	-	See the <i>Cortex-A9 Technical Reference Manual</i> .

Access to coprocessors in the Non-secure state depends on the permissions set in the *Non-secure Access Control Register* on page 2-10.

Attempts to read or write the CPACR Register access bits depend on the corresponding bit for each coprocessor in *Non-secure Access Control Register*. Table 2-4 shows the results of attempted access to coprocessor access bits for each mode.

Table 2-4 Results of access to the CRACR Register

NSACR[11:10]	Secure privileged	Non-secure privileged	Secure or Non-secure User
b00	R/W	RAZ/WI	Access prohibited ^a
b01	R/W	R/W	Access prohibited ^a

a. User privilege access generates an Undefined Instruction exception.

To access the CPACR Register, read or write CP15 with:

```
MRC p15, 0, <Rd>, c1, c0, 2 ; Read Coprocessor Access Control Register
MCR p15, 0, <Rd>, c1, c0, 2 ; Write Coprocessor Access Control Register
```

When the CPACR is updated, the change to the register is guaranteed to be visible only after the next *Instruction Synchronization Barrier (ISB)* instruction. When this register is updated, software must ensure that no instruction that depends on the new or old register values is issued before the ISB instruction.

Normally, software uses a read, modify, write sequence to update the CPACR, to avoid unwanted changes to the access settings for other coprocessors.

———— **Note** ————

You must enable CP10 and CP11 in the CPACR Register before accessing any Advanced SIMD or VFP system registers.

2.5.2 Non-secure Access Control Register

The NSACR Register defines the Non-secure access rights for the Cortex-A9 NEON MPE and other system functionality.

The NSACR Register is:

- a read/write register in Secure state
- a read-only register in Non-secure state
- only accessible in privileged modes.

Figure 2-2 on page 2-11 shows the bit assignments of the NSACR Register relevant to the Cortex-A9 MPE. See the *Cortex-A9 Technical Reference Manual* for details of other fields in this register.

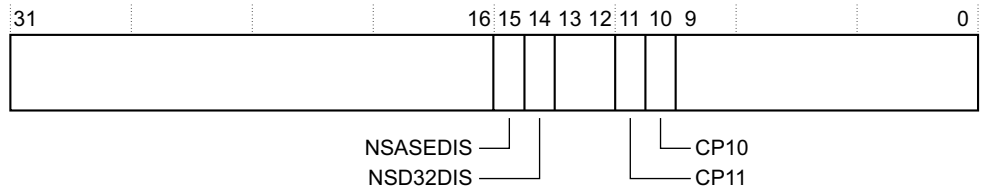


Figure 2-2 NSACR Register bit assignments

Table 2-5 shows the NSACR Register bit assignments.

Table 2-5 NSACR Register bit assignments

Bits	Field	Function
[31:16]	-	See the <i>Cortex-A9 Technical Reference Manual</i> .
[15]	NSASEDIS	Disable Non-secure Advanced SIMD extension functionality: b0 = Full access provided to CPACR.ASEDIS. b1 = The CPACR.ASEDIS bit when executing in Non-secure state has a fixed value of 1 and writes to it are ignored.
[14]	NSD32DIS	Disable Non-secure use of D16-D31 of the VFP register file: b0 = Full access provided to CPACR.D32DIS. b1 = The CPACR.D32DIS bit when executing in Non-secure state has a fixed value of 1 and writes to it are ignored.
[13:12]	-	See the <i>Cortex-A9 Technical Reference Manual</i> .
[11]	CP11	Permission to access coprocessor 11: b0 = Secure access only. This is the reset value. b1 = Secure or Non-secure access.
[10]	CP10	Permission to access coprocessor 10: b0 = Secure access only. This is the reset value. b1 = Secure or Non-secure access.
[9:0]	-	See the <i>Cortex-A9 Technical Reference Manual</i> .

To access the NSACR Register, read or write CP15 with:

MRC p15, 0, <Rd>, c1, c1, 2 ; Read Non-secure Access Control Register data
MCR p15, 0, <Rd>, c1, c1, 2 ; Write Non-secure Access Control Register data

Table 2-6 on page 2-12 shows the results of attempted access for each mode.

Table 2-6 Results of access to the NSACR Register

Secure privileged		Non-secure privileged		User	
Read	Write	Read	Write	Read	Write
Data	Data	Data	Undefined Instruction exception	Undefined Instruction exception	Undefined Instruction exception

2.6 Register summary

Table 2-7 shows the Cortex-A9 NEON MPE system registers. All NEON MPE system registers are 32-bit wide. Reserved register addresses are RAZ/WI.

Table 2-7 Cortex-A9 NEON MPE system registers

Name	Type	Reset	Description
FPSID	RO	0x41033092	See <i>Floating-Point System ID Register</i> on page 2-14
FPSCR	RW	0x00000000	See <i>Floating-Point Status and Control Register</i> on page 2-15
MVFR1	RO	0x01111111	See the <i>ARM Architecture Reference Manual</i>
MVFR0	RO	0x10110222	See the <i>ARM Architecture Reference Manual</i>
FPEXC	RW	0x00000000	See <i>Floating-Point Exception Register</i> on page 2-17

Table 2-8 shows the processor modes for accessing the Cortex-A9 NEON MPE system registers.

Table 2-8 Accessing Cortex-A9 NEON MPE system registers

Register	Privileged access		User access	
	FPEXC EN=0	FPEXC EN=1	FPEXC EN=0	FPEXC EN=1
FPSID	Permitted	Permitted	Not permitted	Not permitted
FPSCR	Not permitted	Permitted	Not permitted	Permitted
MVFR0, MVFR1	Permitted	Permitted	Not permitted	Not permitted
FPEXC	Permitted	Permitted	Not permitted	Not permitted

2.7 Register descriptions

This section describes the Cortex-A9 NEON MPE system registers. Table 2-7 on page 2-13 provides cross references to individual registers.

2.7.1 Floating-Point System ID Register

The FPSID Register characteristics are:

Purpose Provides information about the VFP implementation.

Usage constraints Only accessible in privileged modes.

Configurations Available in all NEON MPE configurations.

Attributes See the register summary in Table 2-7 on page 2-13.

Figure 2-3 shows the FPSID Register bit assignments.

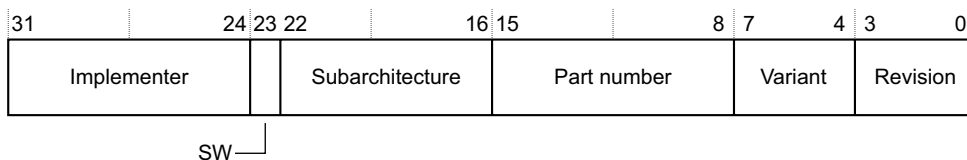


Figure 2-3 FPSID Register bit assignments

Table 2-9 shows the FPSID Register bit assignments.

Table 2-9 FPSID Register bit assignments

Bits	Name	Function
[31:24]	Implementer	Denotes ARM
[23]	SW	Hardware implementation with no software emulation
[22:16]	Subarchitecture	The null VFP sub-architecture
[15:8]	Part number	VFPv3
[7:4]	Variant	Cortex-A9
[3:0]	Revision	Revision 2

You can access the FPSID Register with the following VMRS instruction:

VMRS <Rd>, FPSID ; Read Floating-Point System ID Register

2.7.2 Floating-Point Status and Control Register

The FPSCR characteristics are:

- Purpose** Provides User level control of the FPU.
- Usage constraints** There are no usage constraints.
- Configurations** Available in all FPU configurations.
- Attributes** See the register summary in Table 2-7 on page 2-13.

Figure 2-4 shows the FPSCR bit assignments.

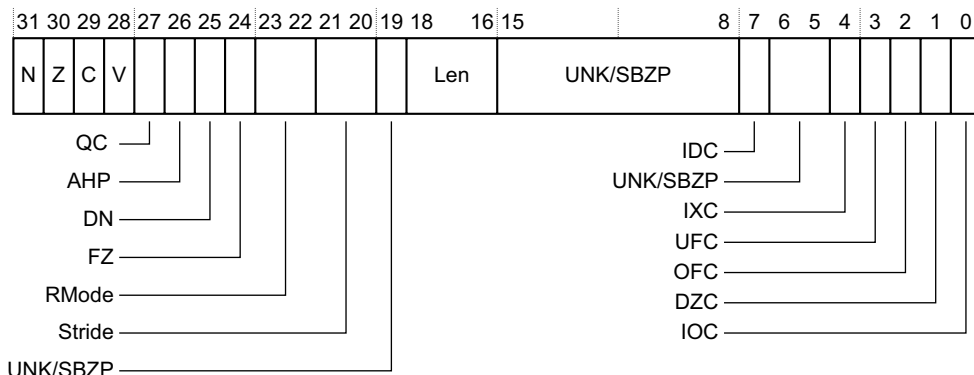


Figure 2-4 FPSCR bit assignments

Table 2-10 shows the FPSCR bit assignments.

Table 2-10 FPSCR bit assignments

Bits	Field	Function
[31]	N	Set to 1 if a comparison operation produces a less than result.
[30]	Z	Set to 1 if a comparison operation produces an equal result.
[29]	C	Set to 1 if a comparison operation produces an equal, greater than, or unordered result.
[28]	V	Set to 1 if a comparison operation produces an unordered result.
[27]	QC	Set to 1 if an Advanced SIMD integer operation has saturated since 0 was last written to this bit. ^a
[26]	AHP	Alternative Half-Precision control bit: b0 = IEEE half-precision format selected b1 = Alternative half-precision format selected.

Table 2-10 FPSCR bit assignments (continued)

Bits	Field	Function
[25]	DN	Default NaN mode control bit: b0 = NaN operands propagate through to the output of a floating-point operation b1 = Any operation involving one or more NaNs returns the Default NaN. Advanced SIMD arithmetic always uses the Default NaN setting, regardless of the value of the DN bit.
[24]	FZ	Flush-to-zero mode control bit: b0 = Flush-to-zero mode disabled. Behavior of the floating-point system is fully compliant with the IEEE 754 standard. b1 = Flush-to-zero mode enabled. Advanced SIMD arithmetic always uses the Flush-to-zero setting, regardless of the value of the FZ bit.
[23:22]	RMode	Rounding Mode control field: b00 = <i>round to nearest</i> (RN) mode b01 = <i>round towards plus infinity</i> (RP) mode b10 = <i>round towards minus infinity</i> (RM) mode b11 = <i>round towards zero</i> (RZ) mode. Advanced SIMD arithmetic always uses the Round to Nearest setting, regardless of the value of the RMode bits.
[21:20]	Stride	Stride control used for backwards compatibility with short vector operations. The Cortex-A9 NEON MPE ignores the value of this field. See the <i>ARM Architecture Reference Manual</i> .
[19]	-	UNK/SBZP.
[18:16]	Len	Vector length, used for backwards compatibility with short vector operation. If you set this field to a non-zero value, the VFP data-processing instructions generate exceptions. See the <i>ARM Architecture Reference Manual</i> .
[15:8]	-	UNK/SBZP.
[7]	IDC	Input Denormal cumulative exception flag. ^a
[6:5]	-	UNK/SBZP.
[4]	IXC	Inexact cumulative exception flag. ^a
[3]	UFC	Underflow cumulative exception flag. ^a

Table 2-10 FPSCR bit assignments (continued)

Bits	Field	Function
[2]	OFC	Overflow cumulative exception flag. ^a
[1]	DZC	Division by Zero cumulative exception flag. ^a
[0]	IOC	Invalid Operation cumulative exception flag. ^a

a. The exception flags, bit [27], bit [7], and bits [4:0] of the FPSCR are exported on the **DEFLAGS** output so they can be monitored externally to the processor, if required.

You can access the FPSCR with the following VMSR instructions:

VMRS <Rd>, FPSCR ; Read Floating-Point Status and Control Register
 VMSR FPSCR, <Rt> ; Write Floating-Point Status and Control Register

2.7.3 Floating-Point Exception Register

The FPEXC Register characteristics are:

- Purpose** Provides global enable control of the Advanced SIMD and VFP extensions.
- Usage constraints**
 - Only accessible in the Non-secure state if the CP10 and CP11 bits in the NSACR are set to 1, see *Non-secure Access Control Register* on page 2-10.
 - Only accessible in privileged modes, and only if access to coprocessors CP10 and CP11 is enabled in the Coprocessor Access Control Register, see *Coprocessor Access Control Register* on page 2-8.
- Configurations** Available in all NEON MPE configurations.
- Attributes** See the register summary in Table 2-7 on page 2-13.

Figure 2-5 on page 2-18 shows the FPEXC Register bit assignments.

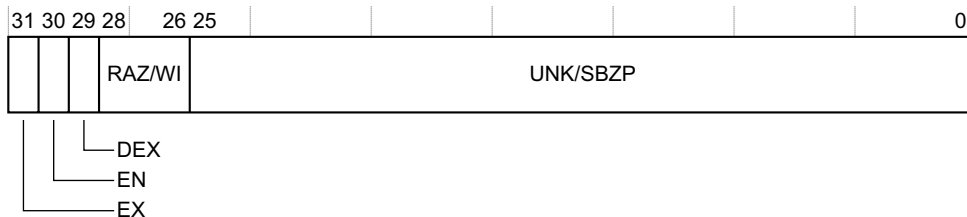


Figure 2-5 FPEXC Register bit assignments

Table 2-11 shows the FPEXC Register bit assignments.

Table 2-11 FPEXC Register bit assignments

Bits	Name	Function
[31]	EX	Exception bit: This bit reads-as-zero and ignores-writes. The Cortex-A9 NEON MPE never requires asynchronous exception handling.
[30]	EN	Enable bit: b0 = the Advanced SIMD and VFP extensions are disabled b1 = the Advanced SIMD and VFP extensions are enabled and operate normally. The EN bit is cleared to 0 at reset.
[29]	DEX	Defined synchronous instruction exceptional flag: b0 = no exception has occurred b1 = attempt to perform a VFP vector operation has been trapped ^a . The DEX bit is cleared to 0 at reset.
[28:26]	-	RAZ/WI.
[25:0]	-	UNK/SBZP.

- a. The Cortex-A9 NEON MPE hardware does not support the deprecated VFP short vector feature. Attempts to execute VFP data-processing instructions when the FPSCR.LEN field is non-zero result in the FPSCR.DEX bit being set and a synchronous Undefined instruction exception being taken. You can use software to emulate the short vector feature, if required.

You can access the FPEXC Register with the following VMRS instructions:

VMRS <Rd>, FPEXC ; Read Floating-Point Status and Control Register
 VMRS FPEXC, <Rt> ; Write Floating-Point Status and Control Register

Chapter 3

Instruction Timing

This chapter describes the cycle timings of instructions on the Cortex-A9 NEON MPE. It contains the following sections:

- *About instruction cycle timing* on page 3-2
- *Writing optimal VFP and Advanced SIMD code* on page 3-3
- *Cortex-A9 NEON MPE instructions* on page 3-4
- *Instruction-specific scheduling* on page 3-10.

3.1 About instruction cycle timing

This chapter provides information to estimate how much execution time particular code sequences require. The complexity of the Cortex-A9 processor makes it impossible to guarantee precise timing information with hand calculations. The timing of an instruction is often affected by other concurrent instructions, memory system activity, and events outside the instruction flow.

Detailed descriptions of all possible instruction interactions and events taking place in the processor are beyond the scope of this document.

3.2 Writing optimal VFP and Advanced SIMD code

The following guidelines can provide significant performance increases for VFP and Advanced SIMD code:

Where possible avoid:

- unnecessary accesses to the VFP control registers
- transferring values between the Cortex-A9 core registers and VFP or Advanced SIMD register file, see the *ARM Architecture Reference Manual* for definition of core registers
- register dependencies between neighboring instructions
- mixing Advanced SIMD only instructions with VFP only instructions.

Be aware that:

- with the exception of simultaneous loads and stores, the processor can execute VFP and Advanced SIMD instructions in parallel with ARM or Thumb instructions
- using Advanced SIMD value selection operations is more efficient than using the equivalent VFP compare with conditional execution.

3.3 Cortex-A9 NEON MPE instructions

Table 3-1 shows the instructions supported by the Cortex-A9 NEON MPE, and the instruction set that they are in, either Advanced SIMD or VFP. The data types that each instruction supports are as follows:

D	Double precision floating-point values
F	Single precision floating-point values
H	Half precision floating-point values
I	Integer values
P	Polynomials with single-bit coefficients
X	Operation is independent of data representation.

See the *ARM Architecture Reference Manual* for details of instruction encodings and functionality.

Table 3-1 Cortex-A9 MPE instructions

Name	Advanced SIMD	VFP	Description
VABA	I	-	Absolute Difference and Accumulate
VABAL	I	-	Absolute Difference and Accumulate Long
VABD	I, F	-	Absolute Difference
VABDL	I	-	Absolute Difference Long
VABS	I, F	F, D	Absolute
VACGE	F	-	Absolute Compare Greater Than or Equal
VACGT	F	-	Absolute Compare Greater Than
VACLE	F	-	Absolute Compare Less Than or Equal
VACLT	F	-	Absolute Compare Less Than
VADD	I, F	F, D	Add
VADDHN	I	-	Add and Narrow Returning High Half
VADDL	I	-	Add Long
VADDW	I	-	Add Wide
VAND	X	-	Bitwise AND
VBIC	I	-	Bitwise Clear

Table 3-1 Cortex-A9 MPE instructions (continued)

Name	Advanced SIMD	VFP	Description
VBIF	X	-	Bitwise Insert if False
VBIT	X	-	Bitwise Insert if True
VBSL	X	-	Bitwise Select
VCEQ	I, F	-	Compare Equal
VCGE	I, F	-	Compare Greater Than or Equal
VCLE	I, F	-	Compare Less Than or Equal
VCLS	I	-	Count Leading Sign Bits
VCLT	I, F	-	Compare Less Than
VCLZ	I	-	Count Leading Zeros
VCMP	-	F, D	Compare Setting Flags
VCNT	I	-	Count Number of Set Bits
VCVT	I, F, H	I, F, D, H	Convert Between Floating-Point and 32-bit Integer Types
VDIV	-	F, D	Divide
VDUP	X	-	Duplicate
VEOR	X	-	Bitwise Exclusive OR
VEXT	I	-	Extract Elements and Concatenate
VHADD	I	-	Halving Add
VHSUB	I	-	Halving Subtract
VLD1	X	-	Load Single-Element Structures
VLD2	X	-	Load Two-Element Structures
VLD3	X	-	Load Three-Element Structures
VLD4	X	-	Load Four-Element Structures
VLDM	X	F, D	Load Multiple Registers
VLDR	X	F, D	Load Single Register

Table 3-1 Cortex-A9 MPE instructions (continued)

Name	Advanced SIMD	VFP	Description
VMAX	I, F	-	Maximum
VMIN	I, F	-	Minimum
VMLA	I, F	F, D	Multiply Accumulate
VMLS	I, F	F, D	Multiply Subtract
VMLAL	I	-	Multiply Accumulate Long
VMLSL	I	-	Multiply Subtract Long
VMOV	X	F, D	Move Register or Immediate
VMOVL	I	-	Move Long
VMOVN	I	-	Move and Narrow
VMRS	X	F, D	Move Advanced SIMD or VFP Register to ARM Compute Engine
VMSR	X	F, D	Move ARM Core Register to Advanced SIMD or VFP
VMUL	I, F, P	F, D	Multiply
VMULL	I, F, P	-	Multiply Long
VMVN	X	-	Bitwise NOT
VNEG	I, F	F, D	Negate
VNMLA	-	F, D	Negative Multiply Accumulate
VNMLS	-	F, D	Negative Multiply Subtract
VNMUL	-	F, D	Negative Multiply
VORN	X	-	Bitwise OR NOT
VORR	X	-	Bitwise OR
VPADAL	I	-	Pairwise Add and Accumulate Long
VPADD	I, F	-	Pairwise Add
VPADDL	I	-	Pairwise Add Long
VPMAX	I, F	-	Pairwise Maximum

Table 3-1 Cortex-A9 MPE instructions (continued)

Name	Advanced SIMD	VFP	Description
VPMIN	I, F	-	Pairwise Minimum
VPOP	X	F, D	Pop from Stack
VPUSH	X	F, D	Push to Stack
VQABS	I	-	Saturating Absolute
VQADD	I	-	Saturating Add
VQDMLAL	I	-	Saturating Double Multiply Accumulate Long
VQDMLSL	I	-	Saturating Double Multiply Subtract Long
VQDMULH	I	-	Saturating Doubling Multiply Returning High Half
VQDMULL	I	-	Saturating Doubling Multiply Long
VQMOVN	I	-	Saturating Move and Narrow
VQMOVUN	I	-	Saturating Move and Unsigned Narrow
VQNEG	I	-	Saturating Negate
VQRDMULH	I	-	Saturating Rounding Doubling Multiply Returning High Half
VQRSHL	I	-	Saturating Rounding Shift Left
VQRSHRN	I	-	Saturating Rounding Shift Right Narrow
VQRSHRUN	I	-	Saturating Rounding Shift Right Unsigned Narrow
VQSHL	I	-	Saturating Shift Left
VQSHLU	I	-	Saturating Shift Left Unsigned
VQSHRN	I	-	Saturating Shift Right Narrow
VQSHRUN	I	-	Saturating Shift Right Unsigned Narrow
VQSUB	I	-	Saturating Subtract
VRADDHN	I	-	Rounding Add and Narrow Returning High Half
VRECPE	I, F	-	Reciprocal Estimate
VRECPS	F	-	Reciprocal Step

Table 3-1 Cortex-A9 MPE instructions (continued)

Name	Advanced SIMD	VFP	Description
VREV16	X	-	Reverse in Halfwords
VREV32	X	-	Reverse in Words
VREV64	X	-	Reverse in Doublewords
VRHADD	I	-	Rounding Halving Add
VRSHL	I	-	Rounding Shift Left
VRSHR	I	-	Rounding Shift Right
VRSHRN	I	-	Rounding Shift Right Narrow
VRSQRTE	I, F	-	Reciprocal Square Root Estimate
VRSQRTS	F	-	Reciprocal Square Root Step
VRSRA	I	-	Rounding Shift Right and Accumulate
VRSUBHN	I	-	Rounding Subtract and Narrow Returning High Half
VSHL	I	-	Shift Left
VSHLL	I	-	Shift Left Long
VSHR	I	-	Shift Right
VSHRN	I	-	Shift Right Narrow
VSLI	X	-	Shift Left and Insert
VSQRT	-	F, D	Square Root
VSRA	I	-	Shift Right and Accumulate
VSRI	X	-	Shift Right and Insert
VST1	X	-	Store single-element structures
VST2	X	-	Store two-element structures
VST3	X	-	Store three-element structures
VST4	X	-	Store four-element structures
VSTM	X	F, D	Store Multiple Registers

Table 3-1 Cortex-A9 MPE instructions (continued)

Name	Advanced SIMD	VFP	Description
VSTR	X	F, D	Store Register
VSUB	I, F	F, D	Subtract
VSUBHN	I	-	Subtract and Narrow
VSUBL	I	-	Subtract Long
VSUBW	I	-	Subtract Wide
VSWP	I	-	Swap Contents
VTBL	X	-	Table Lookup
VTBX	X	-	Table Extension
VTRN	X	-	Transpose
VTST	I	-	Test Bits
VUZP	X	-	Unzip
VZIP	X	-	Zip

3.4 Instruction-specific scheduling

Complex instruction dependencies and memory system interactions make it impossible to concisely specify the exact cycle timing of all instructions in all circumstances. The following sections provide timing information that is accurate in most cases. For precise timing, you must use a cycle-accurate model of both your system and the processor.

3.4.1 Instruction timing tables

The tables in this section provide useful timing information for each category of instruction. The definition of each column is as follows:

Instruction	Instruction mnemonics are in the ARM UAL format. For legacy mnemonic equivalents, see the <i>ARM Architecture Reference Manual</i> .
Format	This is the register format of the instruction for the timing information that is provided. Where multiple formats are given in a single row, the timing information applies to all those formats. In some cases, the table provides operand type information when different operands require different computation time. For example, Table 3-2 on page 3-11, the VFP timing table, contains .F and .D UAL modifiers for single and double-precision operations respectively. Where data-type specifiers are omitted, then you can assume the information applies to all specifiers relevant to the particular entry.
Cycles	This is the number of issue cycles the particular instruction consumes, and is the absolute minimum number of cycles per instruction if no operand interlocks are present.
Source	The Source field indicates the execution cycle where the source operands must be available if the instruction is to be allowed to issue. The comma separated list matches that of the Format field, indicating the register that each value applies to. Where two lines are provided for a single format containing quad (Q) registers, this indicates that the source registers are handled independently between top and bottom half double (D) registers. The first line provides the information for the lower half of the quad register and the second line provides the same information for the upper half of the quad register.
Result	The Result field indicates the execution cycle when the result of the operation is ready. At this point, the result might be ready as source operands for consumption by other instructions using forwarding paths. However, some pairs of instructions might have to wait until the value is written back to the register file.

Writeback The Writeback field indicates the execution cycle that the result is committed to the register file. From this cycle, the result of the instruction is available to all other instructions as a source operand.

———— **Note** —————

An instruction might not issue if its writeback to a particular register might occur before that of an earlier instruction to the same register.

This section contains the following:

- *VFP instruction timing*
- *Advanced SIMD integer arithmetic instructions* on page 3-14
- *Advanced SIMD integer multiply instructions* on page 3-16
- *Advanced SIMD integer shift instructions* on page 3-18
- *Advanced SIMD permute instructions* on page 3-19
- *Advanced SIMD floating-point instructions* on page 3-20
- *Advanced SIMD load-store instructions* on page 3-22.

3.4.2 VFP instruction timing

Table 3-2 shows the VFP instruction timing.

Table 3-2 VFP instruction timing

Name	Format	Cycles	Source	Result	Writeback
VADD	.F Sd,Sn,Sm	1	-,1,1	4	4
VSUB	.D Dd,Dn,Dm				
VCVT	.F Sd,Sn .D Dd,Sn	1	-,1	4	4
VMUL	.F Sd,Sn,Sm	1	-,1,1	5	5
VNMUL	.D Dd,Dn,Dm	2	-,1,1	6	6
VMLA ^a	.F Sd,Sn,Sm	1	-,1,1	8	8
VMLS	.D Dd,Dn,Dm	2	-,1,1	9	9
VNMLS					
VNMLA					
VABS	.F Sd,Sn	1	-,1	1	2
VNEG	.D Dd,Dn				

Table 3-2 VFP instruction timing (continued)

Name	Format	Cycles	Source	Result	Writeback
VMOV	Rt, Sn	1	-,1	-	-
	Rt, Rt2, Dn	1	-, -,1	-	-
	Dd, Rt, Rt2	1	-,1,1	1	2
	Sd, RtSd, SnDd, Dn	1	-,1	1	2
	.F Sd, #imm .D Dd, #imm	1	-, -	1	2
VMRS	Rt, FPSCR	1	-,1	-	-
VDIV	.F Sd, Sn, Sm	10	-,1,1	15	15
	.D Dd, Dn, Dm	20	-,1,1	25	25
VSQRT	.F Sd, Sn, Sm	13	-,1,1	17	17
	.D Dd, Dn, Dm	28	-,1,1	32	32
VCMP	.F Sn, Sm	1	1,1	1	4
	.D Dn, Dm				

- a. If a multiply-accumulate follows a multiply or another multiply-accumulate, and depends on the result of that first instruction, then if the dependency between both instructions are of the same type and size, the processor uses a special multiplier accumulator forwarding. This special forwarding means the multiply instructions can issue back-to-back because the result of the first instruction in cycle 5 is forwarded to the accumulator of the second instruction in cycle 4. If the size and type of the instructions do not match, then Dd or Qd is required in cycle 3. This applies to combinations of the multiply-accumulate instructions VMLA, VMLS, VQDMLA, and VQDMLS, and the multiply instructions VMUL and VQDMUL.

3.4.3 VFP load and store instruction timing

Table 3-3 on page 3-13 shows the VFP load/store instruction timing.

The SP registers column gives the number of single-precision registers that are operated on. For load and store multiple operations, the number of registers in the list is equal to either $2p$ or $2p+1$ for an even or odd number of registers respectively. Load and stores operations using addresses that are not 64-bit aligned can incur an additional cycle.

Table 3-3 VFP load and store instruction timing

Name	Format	64-bit aligned	SP registers	Cycles	Source ^a	Result	Writeback
VLDR	.F Sd, []	-	1	1	-	1	1
VLDR	.D Dd, []	Yes	2	1	-	1	1
		No	2	2	-	1,2	1,2
VLDM	.F Rn, {...}	Yes	$2p$	p	-	1, 2, ..., p	1, 2, ..., p
			$2p+1$	$p+1$	-	1, 2, ..., $p+1$	1, 2, ..., $p+1$
		No	$2p$	$p+1$	-	1, 2, ..., $p+1$	1, 2, ..., $p+1$
			$2p+1$	$p+1$	-	1, 2, ..., $p+1$	1, 2, ..., $p+1$
VLDM	.D Rn, {...}	Yes	$2p$	p	-	1, 2, ..., p	1, 2, ..., p
		No	$2p$	$p+1$	-	1, 2, ..., $p+1$	1, 2, ..., $p+1$
VSTR	.F Sd, []	-	1	1	-	-	-
VSTR	.D Dd, []	Yes	2	1	-	-	-
		No	2	2	-	-	-
VSTM	.F Rn, {...}	Yes	$2p$	p	-	-	-
			$2p+1$	$p+1$	-	-	-
		No	$2p$	$p+1$	-	-	-
			$2p+1$	$p+1$	-	-	-
VSTM	.D Rn, {...}	Yes	$2p$	p	-	-	-
		No	$2p$	$p+1$	-	-	-

a. Store instructions can be issued before their operands are calculated.

3.4.4 Advanced SIMD integer arithmetic instructions

Table 3-4 shows the Advanced SIMD integer arithmetic instruction timing.

Table 3-4 Advanced SIMD integer arithmetic instruction timing

Name	Format	Cycles	Source	Result	Writeback
VADD	Dd, Dn, Dm	1	-,2,2	3	6
VAND	Qd, Qn, Qm				
VORR					
VEOR					
VBIC					
VORN					
VSUB	Dd, Dn, Dm Qd, Qn, Qm	1	-,2,1	3	6
VADDL	Qd, Dn, Dm	1	-,1,1	3	6
VSUBL					
VADDW	Qd, Qn, Dm	1	-,2,1	3	6
VSUBW					
VHADD	Dd, Dn, Dm	1	-,2,2	4	6
VRHADD	Qd, Qn, Qm				
VQADD					
VTST					
VADH	Dd, Qn, Qm	1	-,2,2	4	6
VRADH					
VSBH	Dd, Qn, Qm	1	-,2,1	4	6
VRSBH					
VHSUB	Dd, Dn, Dm	1	-,2,1	4	6
VQSUB	Qd, Qn, Qm				
VABD					
VCEQ					
VCGE					
VCGT					
VMAX					
VMIN					

Table 3-4 Advanced SIMD integer arithmetic instruction timing (continued)

Name	Format	Cycles	Source	Result	Writeback
VPMAX VPMIN	Dd, Dn, Dm	1	-,2,1	4	6
VNEG	Dd, Dm Qd, Qm	1	-,1	3	6
VQNEG VQABS	Dd, Dm Qd, Qm	1	-,1	4	6
VABDL	Qd, Dn, Dm	1	-,2,1	4	6
VABS	Dd, Dm Qd, Qm	1	-,2	4	6
VCEQ VCGE VCGT VCLE VCLT	Dd, Dm, #imm Qd, Qm, #imm	1	-,2,-	4	6
VPADD	Dd, Dn, Dm	1	-,1,1	3	6
VPADDL	Dd, Dn Qd, Qn	1	-,1	3	6
VMVN	Dd, DmQd, Qm	1	-,2	3	6
VCLS VCLZ VCNT	Dd, Dm Qd, Qm	1 2	-,2 -,2 -,3	3 3 4	6 6 7
VMOV VMVN	Dd, #imm Qd, #imm	1	-, -	3	6
VORR VBIC	Dd, #imm Qd, #imm	1	2, -	3	6
VBIT VBIF VBSL	Dd, Dn, Dm Qd, Qn, Qm	1 2	2,2,2 2,2,2 3,3,3	3 3 4	6 6 7

Table 3-4 Advanced SIMD integer arithmetic instruction timing (continued)

Name	Format	Cycles	Source	Result	Writeback
VABA	Dd, Dn, Dm	1	3,2,1	6	6
	Qd, Qn, Qm	2	3,2,1 4,3,2	6	6
VABAL	Qd, Dn, Dm	1	3,2,1	6	6
VPADAL	Dd, Dm Qd, Qm	1	3,1	6	6

3.4.5 Advanced SIMD integer multiply instructions

Table 3-5 shows the operation of the Advanced SIMD integer multiply instruction timing.

Table 3-5 Advanced SIMD integer multiply instructions

Name	Format	Cycles	Source	Result	Writeback
VMUL	.8 Dd, Dn, Dm	1	-,2,2	6	6
VQDMLH	.16 Dd, Dn, Dm				
VQRDMLH	.32 Dd, Dn, Dm	2	-,2,1	7	7
	.32 Qd, Qn, Qm	4	-,2,1 -,4,3	7 9	7 9
VMULL	.8 Qd, Dn, Dm	1	-,2,2	6	6
VQDMULL	.16 Qd, Dn, Dm				
	.32 Qd, Dn, Dm	2	-,2,1	7	7
VMLA	.8 Dd, Dn, Dm	1	3,2,2	6	6
VMLS	.16 Dd, Dn, Dm				
	.8 Qd, Qn, Qm	2	3,2,2	6	6
	.16 Qd, Qn, Qm		4,3,3	7	7
	.32 Dd, Dn, Dm	2	3,2,1	7	7
	.32 Qd, Qn, Qm	4	3,2,1 5,4,3	7 9	7 9

Table 3-5 Advanced SIMD integer multiply instructions (continued)

Name	Format	Cycles	Source	Result	Writeback
VMLAL	.8 Qd, Dn, Dm	1	3,2,2	6	6
VMLS	.16 Qd, Dn, Dm				
VQDMLAL	.32 Qd, Dn, Dm	2	3,2,1	7	7
VQDMLS					
VMUL	.16 Dd, Dn, Dm[x]	1	-,2,1	6	6
VQDMLH	.16 Qd, Qn, Dm[x]	2	-,2,1	6	6
VQRDMLH			-,3,1	7	7
	.32 Dd, Dn, Dm[x]	2	-,2,1	7	7
	.32 Qd, Qn, Qm[x]	4	-,2,1	7	7
			-,4,1	9	9
VMULL	.16 Qd, Dn, Dm[x]	1	-,2,1	6	6
VQDMULL	.32 Qd, Dn, Dm[x]	2	-,2,1	7	7
VMLA	.16 Dd, Dn, Dm[x]	1	3,2,1	6	6
VMLS	.16 Qd, Qn, Dm[x]	2	3,2,1	6	6
	.32 Dd, Dn, Dm[x]	2	3,2,1	7	7
	.32 Qd, Qn, Dm[x]	4	3,2,1	7	7
			5,2,1	9	9
VMLAL	.16 Qd, Dn, Dm[x]	1	3,2,1	6	6
VMLS	.32 Qd, Dn, Dm[x]	2	3,2,1	7	7
VQDMLAL					
VQDMLS					

3.4.6 Advanced SIMD integer shift instructions

Table 3-6 shows the Advanced SIMD integer shift instruction timing.

Table 3-6 Advanced SIMD integer shift instruction timing

Name	Format	Cycles	Source	Result	Writeback
VSHR	Dd,Dm,#imm	1	-,1,-	3	6
VSHL	Qd,Qm,#imm				
VQSHL	Dd,Dm,#imm	1	-,1,-	4	6
VRSHR	Qd,Qm,#imm				
VSHRN	Dd,Qm,#imm	1	-,1,-	3	6
VMOVN					
VQSHRN	Dd,Qm,#imm	1	-,1,-	4	6
VQMOVN					
VQSHRN					
VQRSHR					
VSHL	Qd,Dm,#imm	1	-,1,-	3	6
VMOVL	Qd,Dm	1	-,1	3	6
VSLI	Dd,Dm,#imm	1	1,1,-	3	6
VSRI	Qd,Qm,#imm	2	1,1,-	3	6
			2,2,-	4	7
VSHL	Dd,Dm,Dn	1	-,1,1	3	6
	Qd,Qm,Qn	2	-,1,1	3	6
			-,2,2	4	7
VQSHL	Dd,Dm,Dn	1	-,1,1	4	6
VRSHL	Qd,Qm,Qn	2	-,1,1	5	7
VQRSHL			-,2,2		
VSRA	Dd,Dm,#imm	1	3,1,-	4	6
VRSRA	Qd,Qm,#imm	1	3,1,-	6	6

3.4.7 Advanced SIMD permute instructions

Table 3-7 shows the operation of the Advanced SIMD permute instruction timing.

Table 3-7 Advanced SIMD permute instruction timing

Name	Format	Cycles	Source	Result	Writeback
VDUP	Dd, Dm[x]	1	-,1	2	6
VTRN	Qd, Qm[x]				
VSWP	Dd, Dm	1	1,1	2	6
	Qd, Qm	2	1,1 2,2	2 3	6 7
VZIP	Dd, Dm	1	1,1	2	6
	Qd, Qm	3	1,1 2,2	3,4	7,8
VUZP	Dd, Dm	1	1,1	2	6
	Qd, Qm	3	1,2	3	7
VREV	Dd, DmQd, Qm	1	-,1	2	6
VEXT	Dd, Dn, Dm, #imm	1	-,1,1,-	2	6
	Qd, Qn, Qm, #imm	2	-,1,2,-	3	7
VTBL	Dd, {Dn}, Dm	2	-,2,1	3	7
	Dd, {Dn, Dn1}, Dm	2	-,2,2,1	3	7
	Dd, {Dn, Dn1, Dn2}, Dm	3	-,2,2,3,1	4	8
	Dd, {Dn, Dn1, Dn2, Dn3}, Dm	3	-,2,2,3,3,1	4	8
VTBX	Dd, {Dn}, Dm	2	1,2,1	3	7
	Dd, {Dn, Dn1}, Dm	2	1,2,2,1	3	7
	Dd, {Dn, Dn1, Dn2}, Dm	3	1,2,2,3,1	4	8
	Dd, {Dn, Dn1, Dn2, Dn3}, Dm	3	1,2,2,3,3,1	4	8

3.4.8 Advanced SIMD floating-point instructions

Table 3-8 shows the Advanced SIMD floating-point instructions.

Table 3-8 Advanced SIMD floating-point instructions

Name	Format	Cycles	Source	Result	Writeback
VADD	Dd, Dn, Dm	1	-,2,2	5	6
VSUB	Qd, Qn, Qm	2	-,2,2	5	6
VABD			-,3,3	6	7
VMUL					
VCEQ					
VCGE					
VCGT					
VACGE					
VACGT					
VMAX					
VMIN					
VABS	Dd, Dm	1	-,2	5	6
VNEG	Qd, Qm	2	-,2	5	6
VRECPE			-,3	6	7
VRSQRTE					
VCVT					
VCEQ	Dd, Dm, #imm	1	-,2	5	6
VCGE	Qd, Qm, #imm	2	-,2	5	6
VCGT			-,3	6	7
VCLE					
VCLT					
VPADD	Dd, Dn, Dm	1	-,1,1	5	6
VPMAX					
VPMIN					
VMUL	Dd, Dn, Dm[x]	1	-,2,1	5	6
	Qd, Qn, Dm[x]	2	-,2,1	5	6
			-,3,1	6	7

Table 3-8 Advanced SIMD floating-point instructions (continued)

Name	Format	Cycles	Source	Result	Writeback
VMLA	Dd, Dn, Dm	1	3,2,2	9	10
VMLS	Qd, Qn, Qm	2	3,2,2	9	10
			4,3,3	10	11
	Dd, Dn, Dm[x]	1	3,2,1	9	10
	Qd, Qn, Dm[x]	2	3,2,1 4,3,1	9 10	10 11
VRECPS	Dd, Dn, Dm	1	-,2,2	9	10
VRSQRTS	Qd, Qm, Qn	2	-,2,2	9	10
			-,3,3	10	11

3.4.9 Advanced SIMD load-store instructions

Table 3-9 shows the Advanced SIMD load-store instruction timing.

The values in Table 3-9 correspond to the number of issue cycles required within the MPE execution unit. The number of cycles required by the Cortex-A9 integer processor is equal to the number of 64-bit aligned doublewords that the NEON load or store data overlaps.

Table 3-9 Advanced SIMD load-store instructions

Name	Format	Cycles	Source	Result	Writeback
VLD1	{Dd}, []	2	-	2	7
	{Dd}, [@]	1	-	1	6
	{Dd,Dd1}, []	2	-, -	2,2	7,7
	{Dd,Dd1}, [@]	1	-, -	1,1	6,6
	{Dd,Dd1,Dd2}, []	3	-, -, -	2,2,3	7,7,8
	{Dd,Dd1,Dd2}, [@]	2	-, -, -	1,1,2	6,6,7
	{Dd,Dd1,Dd2,Dd3}, []	3	-, -, -, -	2,2,3,3	7,7,8,8
	{Dd,Dd1,Dd2,Dd3}, [@]	2	-, -, -, -	1,1,2,2	6,6,7,7
	{Dd[x]}, []	3	1	4	8
	{Dd[x]}, [@]	2	1	3	7
	{Dd[]}, []	2	-	3	7
	{Dd[]}, [@]	1	-	2	6
	{Dd[], Dd1[]}, []	2	-, -	3,3	7,7
	{Dd[], Dd1[]}, [@]	1	-, -	2,2	6,6

Table 3-9 Advanced SIMD load-store instructions (continued)

Name	Format	Cycles	Source	Result	Writeback
VLD2	{Dd,Dd1}, []	2	-, -	3,3	7,7
	{Dd,Dd1}, [@]	1	-, -	2,2	6,6
	{Dd,Dd1,Dd2,Dd3}, []	3	-, -, -	3,4,3,4	7,8,7,8
	{Dd,Dd1,Dd2,Dd3}, [@]	2	-, -, -	2,3,2,3	6,7,6,7
	{Dd[x],Dd1[x]}, []	3	1,1	4,4	8,8
	{Dd[x],Dd1[x]}, [@]	2	1,1	3,3	7,7
	{Dd[],Dd1[]}, []	2	-, -	3,3	7,7
	{Dd[],Dd1[]}, [@]	1	-, -	2,2	6,6
VLD3	{Dd,Dd1,Dd2}, []	4	-, -, -	4,4,5	8,8,9
	{Dd,Dd1,Dd2}, [@]	3	-, -, -	3,3,4	7,7,8
	{Dd[x],...,Dd2[x]}, []	5	1,1,2	5,5,6	9,9,10
	{Dd[],Dd1[],Dd2[]}, []	3	-, -, -	3,3,4	7,7,8
VLD4	{Dd,Dd1,Dd2,Dd3}, []	4	-, -, -, -	4,4,5,5	8,8,9,9
	{Dd,Dd1,Dd2,Dd3}, [@]	3	-, -, -, -	3,3,4,4	7,7,8,8
	{Dd[x],...,Dd3[x]}, []	5	1,1,2,2	5,5,6,6	9,9,10,10
	{Dd[x],...,Dd3[x]}, [@]	4	1,1,2,2	4,4,5,5	8,8,9,9
	{Dd[],...,Dd3[]}, []	3	-, -, -, -	3,3,4,4	7,7,8,8
	{Dd[],...,Dd3[]}, [@]	2	-, -, -, -	2,2,3,3	6,6,7,7
VST1	{Dd}, []	2	1	-	-
	{Dd}, [@]	1	1	-	-
	{Dd[x]}, []	2	1	-	-
	{Dd[x]}, [@]	1	1	-	-

Table 3-9 Advanced SIMD load-store instructions (continued)

Name	Format	Cycles	Source	Result	Writeback
VST2	{Dd,Dd1}, []	2	1,1	-	-
	{Dd,Dd1}, [@]	1	1,1	-	-
	{Dd[x],Dd1[x]}, []	2	1,1	-	-
	{Dd[x],Dd1[x]}, [@]	1	1,1	-	-
VST3	{Dd,Dd1,Dd2}, []	3	1,1,2	-	-
	{Dd,Dd1,Dd2}, [@]	2	1,1,2	-	-
	{Dd[x],...,Dd2[x]}, []	3	1,1,2	-	-
VST4	{Dd,Dd1,Dd2,Dd3}, []	3	1,1,2,2	-	-
	{Dd,Dd1,Dd2,Dd3}, [@]	2	1,1,2,2	-	-
	{Dd[x],...,Dd3[x]}, []	3	1,1,2,2	-	-
	{Dd[x],...,Dd3[x]}, [@]	2	1,1,2,2	-	-

Appendix A

Revisions

This appendix describes the technical changes between released issues of this book.

Table A-1 Differences between issue A and issue B

Change	Location
Removed text stating that implementation of Security Extensions is optional	Throughout book
Created separate entry for the VSWP instruction	Table 3-7 on page 3-19
Updated values for the VMLA, VMLS, VRECPS, and VRSQRTS instructions	Table 3-8 on page 3-20

Table A-2 Differences between issue B and issue C

Change	Location
Modified the example code for enabling the Advanced SIMD and VFP extensions in ARM Unified Assembly Language	Example 2-1 on page 2-3
Updated FPSCR bit assignments table	Table 2-10 on page 2-15

Table A-3 Differences between issue C and issue D

Change	Location
Value for version updated to 2	Table 2-7 on page 2-13
FPSCR SBZ/WI bits retitled to UNK/SBZP	Table 2-10 on page 2-15 and Figure 2-4 on page 2-15
SBZ/WI bits retitled to UNK/SBZ	Table 2-11 on page 2-18 and Figure 2-5 on page 2-18
Added Footnote about special behavior of two VMLA instructions	Table 3-2 on page 3-11
Footnote reworded	Table 3-3 on page 3-13
Added paragraph about timings	Table 3-9 on page 3-22
Table formatting fixes	Table 3-3 on page 3-13 Table 3-4 on page 3-14 Table 3-5 on page 3-16 Table 3-6 on page 3-18 Table 3-7 on page 3-19 Table 3-8 on page 3-20 Table 3-9 on page 3-22

Table A-4 Differences between issue D and Issue E

Change	Location
No technical change	-

Table A-5 Differences between issue E and Issue F

Change	Location
ISB added to code example	Example 2-1 on page 2-3
Harmonized FPEXC bit descriptions with the FPU TRM descriptions	Table 2-11 on page 2-18

Glossary

This glossary describes some of the terms used in technical documents from ARM.

Abort	A mechanism that indicates to a core that the value associated with a memory access is invalid. An abort can be caused by the external or internal memory system as a result of attempting to access invalid instruction or data memory. An abort is classified as either a Prefetch or Data Abort, and an internal or External Abort.
Addressing modes	A mechanism, shared by many different instructions, for generating values used by the instructions. For four of the ARM addressing modes, the values generated are memory addresses (which is the traditional role of an addressing mode). A fifth addressing mode generates values to be used as operands by data-processing instructions.
Aligned	A data item stored at an address that is divisible by the number of bytes that defines the data size is said to be aligned. Aligned words and halfwords have addresses that are divisible by four and two respectively. The terms word-aligned and halfword-aligned therefore stipulate addresses that are divisible by four and two respectively.
Architecture	The organization of hardware and/or software that characterizes a processor and its attached components, and enables devices with similar characteristics to be grouped together when describing their behavior, for example, Harvard architecture, instruction set architecture, ARMv6 architecture.
ARM instruction	A word that specifies an operation for an ARM processor to perform. ARM instructions must be word-aligned.

ARM state	A processor that is executing ARM (32-bit) word-aligned instructions is operating in ARM state.
Base register	A register specified by a load or store instruction that is used to hold the base value for the instruction's address calculation. Depending on the instruction and its addressing mode, an offset can be added to or subtracted from the base register value to form the virtual address that is sent to memory.
Bounce	The VFP coprocessor bounces an instruction when it fails to signal the acceptance of a valid VFP instruction to the ARM processor. This action initiates Undefined instruction processing by the ARM processor. The VFP support code is called to complete the instruction that was found to be exceptional or unsupported by the VFP coprocessor. <i>See also</i> Trigger instruction, Potentially exceptional instruction, and Exceptional state.
DP instruction	Coprocessor data processing instruction. For the VFP11 coprocessor, CDP instructions are arithmetic instructions and FCPY, FABS, and FNEG. <i>See also</i> Arithmetic instruction.
Condition field	A four-bit field in an instruction that specifies a condition under which the instruction can execute.
Conditional execution	If the condition code flags indicate that the corresponding condition is true when the instruction starts executing, it executes normally. Otherwise, the instruction does nothing.
Control bits	The bottom eight bits of a Program Status Register. The control bits change when an exception arises and can be altered by software only when the processor is in a privileged mode.
Coprocessor	A processor that supplements the main processor. It carries out additional functions that the main processor cannot perform. Usually used for floating-point math calculations, signal processing, or memory management.
CoreSight	The infrastructure for monitoring, tracing, and debugging a complete system on chip.
Default NaN mode	A mode in which all operations that result in a NaN return the default NaN, regardless of the cause of the NaN result. This mode is compliant with the IEEE 754 standard but implies that all information contained in any input NaNs to an operation is lost.
Denormalized value	<i>See</i> Subnormal value.
Disabled exception	An exception is disabled when its exception enable bit in the FPCSR is not set. For these exceptions, the IEEE 754 standard defines the result to be returned. An operation that generates an exception condition can bounce to the support code to produce the result defined by the IEEE 754 standard. The exception is not reported to the user trap handler.

Double-precision value	Consists of two 32-bit words that must appear consecutively in memory and must both be word-aligned, and that is interpreted as a basic double-precision floating-point number according to the IEEE 754-1985 standard.
Doubleword	A 64-bit data item. The contents are taken as being an unsigned integer unless otherwise stated.
Doubleword-aligned	A data item having a memory address that is divisible by eight.
Enabled exception	An exception is enabled when its exception enable bit in the FPCSR is set. When an enabled exception occurs, a trap to the user handler is taken. An operation that generates an exception condition might bounce to the support code to produce the result defined by the IEEE 754 standard. The exception is then reported to the user trap handler.
Exception	A fault or error event that is considered serious enough to require that program execution is interrupted. Examples include attempting to perform an invalid memory access, external interrupts, and undefined instructions. When an exception occurs, normal program flow is interrupted and execution is resumed at the corresponding exception vector. This contains the first instruction of the interrupt handler to deal with the exception.
Exceptional state	<p>When a potentially exceptional instruction is issued, the VFP11 coprocessor sets the EX bit, FPEXC[31], and loads a copy of the potentially exceptional instruction in the FPINST register. If the instruction is a short vector operation, the register fields in FPINST are altered to point to the potentially exceptional iteration. When in the exceptional state, the issue of a trigger instruction to the VFP11 coprocessor causes a bounce.</p> <p><i>See also</i> Bounce, Potentially exceptional instruction, and Trigger instruction.</p>
Exception handling routine	<i>See</i> Interrupt handler.
Flush-to-zero mode	<p>In this mode, the VFP11 coprocessor treats the following values as positive zeros:</p> <ul style="list-style-type: none"> • arithmetic operation inputs that are in the subnormal range for the input precision • arithmetic operation results, other than computed zero results, that are in the subnormal range for the input precision before rounding. <p>The VFP11 coprocessor does not interpret these values as subnormal values or convert them to subnormal values.</p> <p>The subnormal range for the input precision is $-2^{E_{min}} < x < 0$ or $0 < x < 2^{E_{min}}$.</p>
Halfword	A 16-bit data item.

Halfword aligned	<i>See</i> Aligned
IEEE 754 standard	<i>IEEE Standard for Binary Floating-Point Arithmetic, ANSI/IEEE Std. 754-1985.</i> The standard that defines data types, correct operation, exception types and handling, and error bounds for floating-point systems. Most processors are built in compliance with the standard in either hardware or a combination of hardware and software.
Illegal instruction	An instruction that is architecturally Undefined.
Implementation-defined	The behavior is not architecturally defined, but is defined and documented by individual implementations.
Implementation-specific	The behavior is not architecturally defined, and does not have to be documented by individual implementations. Used when there are a number of implementation options available and the option chosen does not affect software compatibility.
Infinity	In the IEEE 754 standard format to represent infinity, the exponent is the maximum for the precision and the fraction is all zeros.
Input exception	A VFP exception condition in which one or more of the operands for a given operation are not supported by the hardware. The operation bounces to support code for processing.
Interrupt handler	A program that control of the processor is passed to when an interrupt occurs.
Load/store architecture	A processor architecture where data-processing operations only operate on register contents, not directly on memory contents.
Memory bank	One of two or more parallel divisions of interleaved memory, usually one word wide, that enable reads and writes of multiple words at a time, rather than single words. All memory banks are addressed simultaneously and a bank enable or chip select signal determines which of the banks is accessed for each transfer. Accesses to sequential word addresses cause accesses to sequential banks. This enables the delays associated with accessing a bank to occur during the access to its adjacent bank, speeding up memory transfers.
NaN	Not a number. A symbolic entity encoded in a floating-point format that has the maximum exponent field and a nonzero fraction. An SNaN causes an invalid operand exception if used as an operand and a most significant fraction bit of zero. A QNaN propagates through almost every arithmetic operation without signaling exceptions and has a most significant fraction bit of one.

Processor	A processor is the circuitry in a computer system required to process data using the computer instructions. It is an abbreviation of microprocessor. A clock source, power supplies, and main memory are also required to create a minimum complete working computer system.
Read	<p>Reads are defined as memory operations that have the semantics of a load. That is, the ARM instructions LDM, LDRD, LDC, LDR, LDRT, LDRSH, LDRH, LDRSB, LDRB, LDRBT, LDREX, RFE, STREX, SWP, and SWPB, and the Thumb instructions LDM, LDR, LDRSH, LDRH, LDRSB, LDRB, and POP.</p> <p>Java bytecodes that are accelerated by hardware can cause a number of reads to occur, according to the state of the Java stack and the implementation of the Java hardware acceleration.</p>
Reserved	A field in a control register or instruction format is reserved if the field is to be defined by the implementation, or produces Unpredictable results if the contents of the field are not zero. These fields are reserved for use in future extensions of the architecture or are implementation-specific. All reserved bits not used by the implementation must be written as 0 and read as 0.
Rounding mode	<p>The IEEE 754 standard requires all calculations to be performed as if to an infinite precision. For example, a multiply of two single-precision values must accurately calculate the significand to twice the number of bits of the significand. To represent this value in the destination precision, rounding of the significand is often required. The IEEE 754 standard specifies four rounding modes.</p> <p>In round-to-nearest mode, the result is rounded at the halfway point, with the tie case rounding up if it would clear the least significant bit of the significand, making it even.</p> <p>Round-towards-zero mode chops any bits to the right of the significand, always rounding down, and is used by the C, C++, and Java languages in integer conversions.</p> <p>Round-towards-plus-infinity mode and round-towards-minus-infinity mode are used in interval arithmetic.</p>
SBO	<i>See Should Be One.</i>
SBZ	<i>See Should Be Zero.</i>
SBZP	<i>See Should Be Zero or Preserved.</i>
Scalar operation	<p>A VFP coprocessor operation involving a single source register and a single destination register.</p> <p><i>See also</i> Vector operation.</p>

Short vector operation

A VFP coprocessor operation involving more than one destination register and perhaps more than one source register in the generation of the result for each destination.

Should Be One (SBO)

Should be written as 1 (or all 1s for bit fields) by software. Writing a 0 produces Unpredictable results.

Should Be Zero (SBZ)

Should be written as 0 (or all 0s for bit fields) by software. Writing a 1 produces Unpredictable results.

Should Be Zero or Preserved (SBZP)

Should be written as 0 (or all 0s for bit fields) by software, or preserved by writing the same value back that has been previously read from the same field on the same processor.

Significand

The component of a binary floating-point number that consists of an explicit or implicit leading bit to the left of the implied binary point and a fraction field to the right.

Stride

The stride field, FPSCR[21:20], specifies the increment applied to register addresses in short vector operations. A stride of 00, specifying an increment of +1, causes a short vector operation to increment each vector register by +1 for each iteration, while a stride of 11 specifies an increment of +2.

Subnormal value

A value in the range $(-2^{E_{\min}} < x < 2^{E_{\min}})$, except for 0. In the IEEE 754 standard format for single-precision and double-precision operands, a subnormal value has a zero exponent and a nonzero fraction field. The IEEE 754 standard requires that the generation and manipulation of subnormal operands be performed with the same precision as normal operands.

Support code

Software that must be used to complement the hardware to provide compatibility with the IEEE 754 standard. The support code has a library of routines that performs supported functions, such as divide with unsupported inputs or inputs that might generate an exception in addition to operations beyond the scope of the hardware. The support code has a set of exception handlers to process exceptional conditions in compliance with the IEEE 754 standard.

Thumb instruction

A halfword that specifies an operation for an ARM processor in Thumb state to perform. Thumb instructions must be halfword-aligned.

Thumb state

A processor that is executing Thumb (16-bit) halfword aligned instructions is operating in Thumb state.

Trap

An exceptional condition in a VFP coprocessor that has the respective exception enable bit set in the FPSCR register. The user trap handler is executed.

Unaligned	A data item stored at an address that is not divisible by the number of bytes that defines the data size is said to be unaligned. For example, a word stored at an address that is not divisible by four.
Undefined	Indicates an instruction that generates an Undefined instruction trap. See the <i>ARM Architecture Reference Manual</i> for more details on ARM exceptions.
UNP	<i>See</i> Unpredictable.
Unpredictable	For reads, the data returned when reading from this location is unpredictable. It can have any value. For writes, writing to this location causes unpredictable behavior, or an unpredictable change in device configuration. Unpredictable instructions must not halt or hang the processor, or any part of the system.
Unsupported values	Specific data values that are not processed by the VFP coprocessor hardware but bounced to the support code for completion. These data can include infinities, NaNs, subnormal values, and zeros. An implementation is free to select which of these values is supported in hardware fully or partially, or requires assistance from support code to complete the operation. Any exception resulting from processing unsupported data is trapped to user code if the corresponding exception enable bit for the exception is set.
Vector operation	A VFP coprocessor operation involving more than one destination register, perhaps involving different source registers in the generation of the result for each destination. <i>See also</i> Scalar operation.
Word	A 32-bit data item.
Word aligned	<i>See</i> Aligned.

